

Regular expressions, Backus Naur and Compilers

Advanced Simulation 2017 IKP - Marius Reed

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Regular expressions

A regular expression is a sequence of characters that defines a search pattern.

- Applications
 - Search
 - Validate
 - Replace
 - Reformat

Single character

Symbol = a

Regex = a

a

A computer once beat me at chess, but it was no match for me at kick boxing. - *Emo Philips*

Dot character (.)

Symbol = .

Regex = a.e

a.e

Software is like entropy: It is difficult to grasp, weighs nothing, and obeys the Second Law of Thermodynamics; i.e., it always increases. - *Norman Augustine*

Ranges

Symbol = []

Regex = `er[ey]`

`er[ey]`

There is only one problem with common sense; it's not very common. - *Milt Bryce*

Shortcuts (ASCII): `[0-9]`, `[a-z]`, `[0-9a-zA-Z]`

Multipliers

- $*$: item occurs zero or more times.
- $+$: item occurs one or more times.
- $?$: item occurs zero or one time.
- $\{a\}$: item occurs exactly a times.
- $\{a,b\}$: item occurs between a and b times.
- $\{a,\}$: item occurs at least a times.

Multipliers

et+

Getting information off the Internet is like taking a drink from a fire hydrant. - *Mitchell Kapor*

Greedy

a.*e

The greatest enemy of knowledge is not ignorance, it is the illusion of knowledge. - *Stephen Hawking*

Not greedy

a.*?e

The greatest enemy of knowledge is not ignorance, it is the illusion of knowledge. - *Stephen Hawking*

Shortcut for character classes

- `\s` : Matches anything which is considered whitespace.
- `\S` : Matches anything which is NOT considered whitespace.
- `\d` : Matches anything considered a digit. (Same as `[0-9]`)
- `\D` : Matches anything which is NOT considered a digit.
- `\w` : Matches anything which is considered a word character. (Same as `[A-Za-z0-9_]`)
- `\W` : Matches anything which is NOT considered a word character.

`\w+`

Imagination is more important than knowledge. For knowledge is limited, whereas imagination embraces the entire world, stimulating progress, giving birth to evolution. - *Albert Einstein*

anchors and word boundaries

- Symbol = `^` and `$`
- Regex = `^\w+` and `\w+ [\.\!\?]? $`

`^\w+`

Your most unhappy customers are your greatest source of learning. - *Bill Gates*

`\w+ [\.\!\?]? $`

Your most unhappy customers are your greatest source of **learning.** - *Bill Gates*

- Symbol = `\b` and `\b`
- Regex = `\bi\w+\b`

`\bi\w+\b`

The best way to get accurate **information** on Usenet **is** to post something wrong and wait for corrections. - *Matthew Austern*

Grouping

- Symbol = ()
- Regex = `(\d{1,3}\.){3}\d{1,3}`

`(\d{1,3}\.){3}\d{1,3}`

This is an IP address **210.1.41.103**, but this is not 112.3332.124.1

`\b((25[0-5]|2[0-4][0-9]|1[0-9][0-9]|[1-9]?[0-9])\.){3}(25[0-5]|2[0-4][0-9]|1[0-9][0-9]|[1-9]?[0-9])\b`

Alternation (or)

- Symbol = |
- Regex = (John|Michael) Smith

(John|Michael) Smith

We want to match John Smith and Michael Smith, but not Paul Smith.

Example: NTNU email address

- Alt. 1 = `[a-z]+@(stud\.)?ntnu.no`
- Alt. 2 = `[a-z]+@(stud\.ntnu\.no|ntnu\.no)`

`[a-z]+@(stud\.)?ntnu.no`

`mariusre@stud.ntnu.no; mariusre@ntnu.no; marius.reed@lyse.net; john.smith@gmail.com`

Re module in python

- `re.compile(pattern, flags=0)`
- `re.search(pattern,string,flags=0)`
- `re.match(pattern,string,flags=0)`
- `re.split(pattern,string,maxsplit=0, flags = 0)`
- `re.findall(pattern,string,flags=0)`
- `re.finditer(pattern,string,flags=0)`
- `re.sub(pattern,repl,string,count=0,flags=0)`

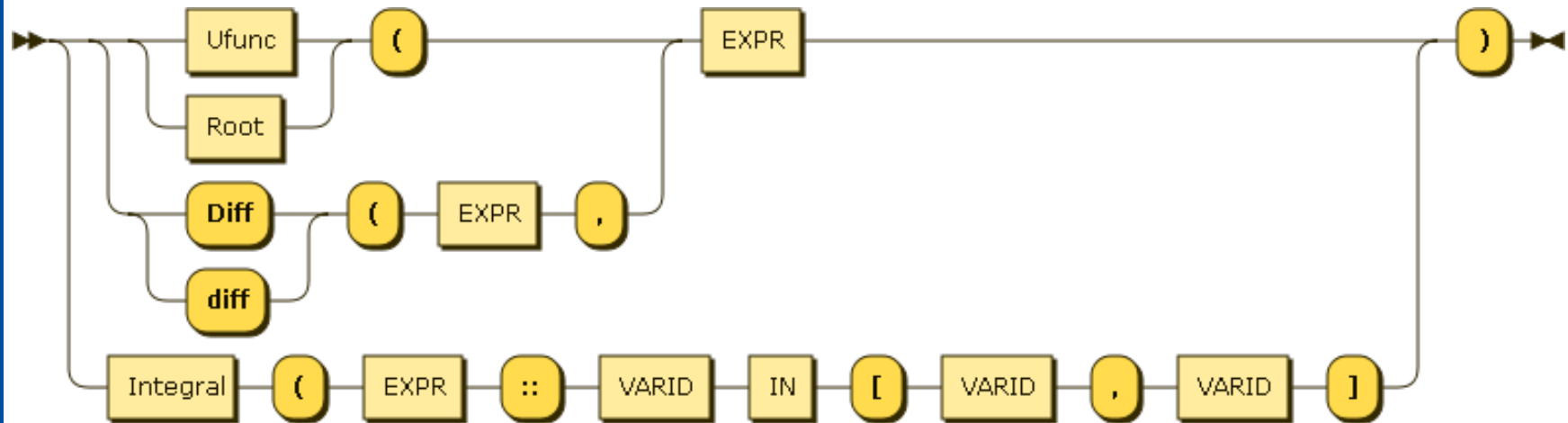
Backus-Naur form

- Notation technique for context free grammar
- Describe syntax

```
<postal-address> ::= <name-part> <street-address> <zip-part>
<name-part>      ::= <personal-part> <last-name> | <personal-part> <name-part>
<personal-part>  ::= <initial> "." | <first-name>
<street-address> ::= <street-name> <house-num> <opt-apt-num>
<zip-part>       ::= <ZIP-code> <town-name>
<opt-apt-num>    ::= <apt-num> | ""
```

Rail road diagram (Syntax Diagram)

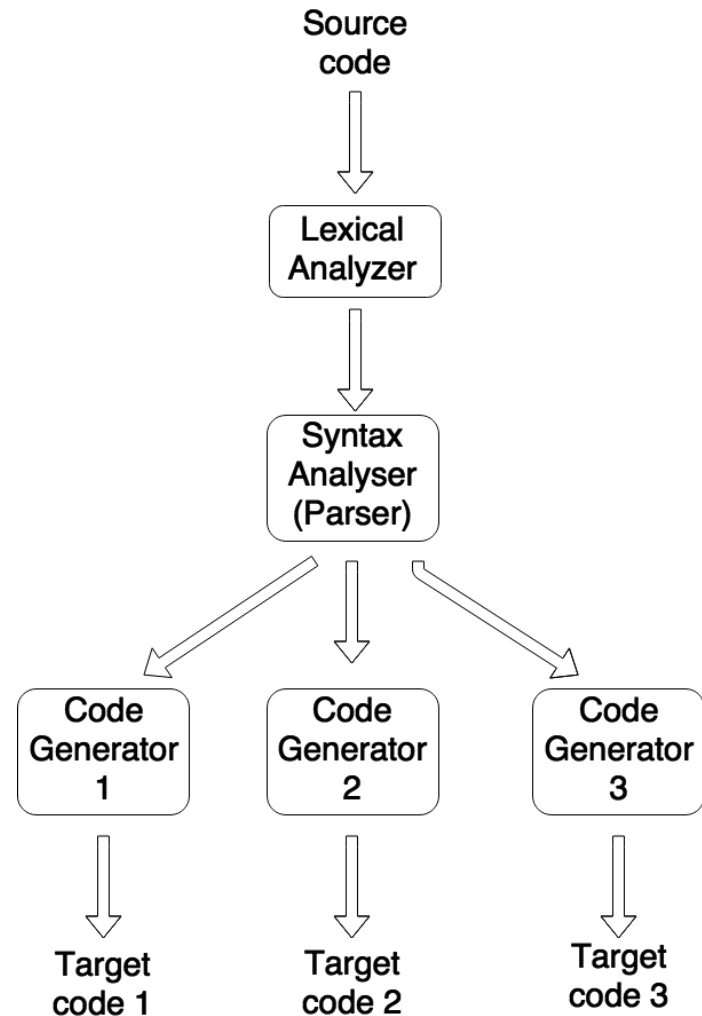
$\langle \text{FUNC} \rangle ::= \langle \text{Ufunc} \rangle "(" \langle \text{EXPR} \rangle ")" \mid \langle \text{Integral} \rangle "(" \langle \text{EXPR} \rangle "::" \langle \text{VARID} \rangle$
 $\langle \text{IN} \rangle "[" \langle \text{VARID} \rangle "," \langle \text{VARID} \rangle "]" "(" \mid \langle \text{Root} \rangle "(" \langle \text{EXPR} \rangle ")" \mid$
 $"\text{Diff}" "(" \langle \text{EXPR} \rangle "," \langle \text{EXPR} \rangle ")" \mid "diff" "(" \langle \text{EXPR} \rangle "," \langle \text{EXPR} \rangle$
 $)"$



Compilers

"A compiler is a special program that processes statements written in a particular programming language and turns them into machine language or "code" that a computer's processor uses"

- Lexical analyzer
- Parser
- Code Generator



Objective

```
[44]
lhs = ndiff
incidence_list = ['Fn_d', 'Fn_v', 'Pnrn', 'fn_d', 'fn_v', 'tnr']
rhs = Fn_v|arc_v & species|fn_v+Fn_d|arc_d & species|fn_d+Pnrn|node_conversion &
  ↪ species|tnr
layer = ontology_physical
```



```
ndiff = reduceProduct(Fn_v,add(fn_v,reduceProduct(Fn_d,add(fn_d,reduceProduct(
  ↪ Pnrn,tnr,node_conversion & species)),arc_d & species)),arc_v &
  ↪ species)
```


Lua and LPeg

- PEG = Parsing expression grammar
 - Choice operator is ordered
- LPeg
 - LPeg patterns
 - Recursive descent parser

LPEG	Description	re python
<code>lpeg.P(<i>string</i>)</code>	Creates a pattern that matches <i>string</i> literally	<code>re.compile(<i>string</i>)</code>
<code>lpeg.S(<i>string</i>)</code>	Creates a pattern that matches any character in <i>string</i>	<code>re.compile([<i>string</i>])</code>
<code>lpeg.R("xy")</code>	Creates a pattern that matches any character between x and y	<code>re.compile([x-y])</code>
<code>patt1 * patt2</code>	Matches patt1 followed by patt2	Not introduced
<code>patt1 + patt2</code>	Matches patt1 or patt2	<code>patt1 patt2</code>
<code>pattⁿ</code>	Matches at least n repetitions of patt	<code>patt{n,}</code>
<code>patt⁻ⁿ</code>	Matches at most n repetitions of patt	<code>patt{,n}</code>

Lexer

```
-- Lexer
spc = L.S(" \t\n")^0
Ufunc      = spc * L.P("inv") + L.P("sqrt") + L.P("exp") + L.P("log") +
  → L.P("sin")
           + L.P("cos") + L.P("tan") + L.P("asin") + L.P("acos")
           + L.P("atan") + L.P("sign") + L.P("neg") + L.P("abs")
           + L.P("min") + L.P("max") * spc
Root       = spc * L.P("root") * spc
Left       = spc * L.P("<") * spc
Right      = spc * L.P(">") * spc
IN         = spc * L.P("in") * spc
Integral   = spc * L.P("integral") * spc
Selection  = spc * L.P("select") * spc
Instantiate = spc * L.P("set") * spc
VarID      = spc * L.R("az", "AZ", "__", "09")^0 * spc
sum        = spc * L.S('+ -') * spc
power      = spc * L.P("^") * spc
dot        = spc * L.P(".") * spc
DL         = spc * L.P("|") * spc
KR         = spc * L.P("'") * spc
par        = spc * L.P("(") + L.P(")") * spc
real       = spc * L.C(
           L.P('-')^1 *
           L.R('09')^0 *
           ( L.P('.') *
             L.R('09')^0
           )^1 ) * spc /
tonumber
integer    = L.R("09")^1
```

Parser

- *The parser will parse the token sequence to identify the syntactic structure of the program.*
- “Connected” to the source language, not the target language.
- *Captures* tokens that is sent to the code generator

Parser and Backus-Naur form

$\langle \text{START} \rangle ::= \langle \text{INstantiate} \rangle "(" \langle \text{ATOM} \rangle ")" | \langle \text{EXPR} \rangle$

$\langle \text{EXPR} \rangle ::= \langle \text{TERM} \rangle \text{ sum } \langle \text{EXPR} \rangle | \langle \text{SEL} \rangle$

$\langle \text{TERM} \rangle ::= "(" \langle \text{EXPR} \rangle ")" | \langle \text{FACT} \rangle \text{ dot } \langle \text{EXPR} \rangle | \langle \text{FACT} \rangle \text{ KR } \langle \text{EXPR} \rangle |$
 $\langle \text{FACT} \rangle \text{ DL } \langle \text{INDID} \rangle \text{ DL } \langle \text{EXPR} \rangle | \langle \text{FACT} \rangle$

```
myGrammar = L.P{
    "START";
    START = (
        L.Ct(Instantiate * "(" * L.V("ATOM") * ")")/instant
        + L.V("EXPR")
    ),
    EXPR = (
        L.Ct(L.V("TERM") * L.C(sum) * L.V("EXPR"))/add
        + L.V("SEL")
    ),
    TERM = (
        L.Ct("(" * L.V("EXPR") * ")") /group
        + L.Ct(L.V("FACT") * dot * L.V("EXPR"))/ExpandProduct
        + L.Ct(L.V("FACT") * KR * L.V("EXPR"))/KhatiriRao
        + L.Ct(L.V("FACT") * DL * L.V("INDID") * DL *
            → L.V("EXPR"))/ReduceProduct
        + L.V("FACT")
    ),
}
```

Expression	Function called
$\underbrace{\text{half.}(\text{abs}(\text{Fm_vv}) + (\text{Fm_vv.D_v}))}_{\text{EXPR}}$	
$\underbrace{\underbrace{\text{half}}_{\text{FACT}} \cdot (\text{abs}(\text{Fm_vv}) + (\text{Fm_vv.D_v}))}_{\text{EXPR}}$	ExpandProduct(FACT, EXPR)
$\underbrace{\text{half}}_{\text{VARID}}$	Just captured
$\underbrace{\text{abs}(\text{Fm_vv})}_{\text{TERM}} \underbrace{+}_{\text{sum}} \underbrace{(\text{Fm_vv.D_v})}_{\text{EXPR}}$	add(TERM, EXPR)
$\underbrace{\text{abs}}_{\text{Ufunc}} (\underbrace{\text{Fm_vv}}_{\text{EXPR}})$	abs(EXPR)
$\underbrace{\text{Fm_vv}}_{\text{VARID}}$	Just captured
$\underbrace{\text{Fm_vv}}_{\text{FACT}} \cdot \underbrace{\text{D_v}}_{\text{EXPR}}$	ExpandProduct(FACT, EXPR)
$\underbrace{\text{Fm_vv}}_{\text{VARID}}$	Just captured
$\underbrace{\text{D_v}}_{\text{VARID}}$	Just captured
TOTAL	ExpandProduct(half, add(abs(Fm_vv),ExpandProduct(Fm_vv, D_v)))

Code generation

- Captures from the parser is sent to the code generator
- In this case, several Lua functions

```
myGrammar = L.P{  
    "START";  
    START = (  
        L.Ct(Instantiate * "(" * L.V("ATOM") * ")")/instant  
        + L.V("EXPR")  
    ),  
}
```

```
function instant(myList)  
    return ("Instantiate(" .. myList[1] .. ")" )  
end
```